

A Lightweight Operating System for the SSDP

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LeanOS – a Lean Operating System

- nationally funded by the FFG under project number 847987
- based upon collaboration between RUAG Space Austria (RSA) and University of Vienna (UVIE) (ESA contract number: 40000107815/13/NL/EL/f)
- tailored to Network-On-Chip/DSP concept of the SSDP
- released under an Open Source license

Background: SSDP/MPPB



- LEON processor (controller)
- Network-on-Chip (NoC)
- 2 Xentium VLIW DSP cores
- High-speed interfaces (SpW, ADC/DAC)
- 50 MHz system clock
- current version 2.0 with SSDP-like feature set and characteristics



Background: SSDP/MPPB



Network-On-Chip

- DMA
 - high-speed packet-like data transfers between devices
 - 2D stride support
 - parallel channels
- Xentium DSP subsystem
 - 10 parallel execution units per DSP core
 - Local tightly-coupled memory (TCM): 4 x 8 kiB banks
 - I-cache: 16 kiB
 - very high performance; example: up to 1.3 s/c computable when sampling up the ramp





Key User Requirements

- it shall be lean and efficient
- it shall support a scalable number of Xentium DSP cores
- it shall be easy to use
- it shall support Fault Detection, Isolation and Recovery (FDIR)
- it shall be designed with applicable S/W standards to be space qualifyable
- it shall come with support documentation and demo applications



Key Functional Requirements

- trap/interrupt handling
- SMP readiness
- threads, scheduling
- MMU support, paging and virtual memory
- loadable module support
- run-time configuration interface
- SSDP hardware drivers (NoC DMA, Xentium, I/O ...)





- architecture-dependent code
 - interfaces to underlying hardware architecture
 - implements abstractions to higher-level kernel code
- kernel code
 - implements services and drivers
 - components shared between architectures
- system call interface
 - provides call interface to user space





Features: Memory Management

- dynamic memory (de-)allocation at run-time
- SRMMU support configurable
 - deferred page allocation via sbrk()
 - virtual memory, 4k paging
- systems without MMU supported, uses same underlying allocator



Features: Loadable Modules

- load/unload kernel modules/drivers at run-time
- easy patching
- easy reconfiguration depending on operating mode
- compiled as relocatable object codes
- built modules are added to payload image
- symbols and dependencies are resolved on load
- may be compiled-in



Features: Configuration/Building

- derived from Linux kbuild system
- easy to use
- prevents configuration conflicts
- also builds Xentium programs
- creates executable and payload image

config - LeanOS Configuration
LeanOS Configuration Arrow keys navigate the menu. <enter> selects submenus> (or empty submenus). Highlighted letters are hotkeys. Pressing <y> includes, <n> excludes, <m> modularizes features. Press <esc><esc> to exit, <? > for Help, for Search.</esc></esc></m></n></y></enter>
<pre>SPARC Configuration> General Setup> Enable loadable module support I Build Xentium DSP kernel programs Compile-time checks and compiler options> [*] Enable printing in kernel (7) Kernel debug level I Enable IRQ statistics via sysctl I Enable System-On-Chip configurations Core Components> I Sample and demo code</pre>
<pre></pre>

Xentium Programs



Design Decisions

- conventional approach:
 - single monolithic program running on DSP
 - data processing is implemented as a series of function calls/operations
- downsides:
 - code size may exceed Xentium i-cache (16 kiB), requiring costly re-fetches via NoC
 - stack is in TCM (4x8 kiB), excessive call depths may use up valuable memory
 - even small changes require full program re-validation

Xentium Programs



Design Decisions

- alternative approach developed by UVIE:
 - multiple tiny programs, one for each operational step (≈ function call)
 - programs can (and should) be dumb, i.e. perform only one task, regardless of larger application
- upsides:
 - code can be arbitrarily small, thus will always fit the i-cache (break up into sub-components)
 - processing chains can be created from independent building blocks
 - changes affect isolated components only (delta-testing)
 - simpler units generally have less execution paths and are easier to trace for WCET

nodes are similar to function calls in a program

processing nodes are executed as needed

Key Feature: Xentium Processing Network

- metadata task packets attach to nodes
- metadata describes data product

Xentium Programs

- "fingerprints" of completed and pending processing operations
- task packets are automatically routed to next matching processing node
- different types of products create their pipeline on-the-fly



Processing Task



Xentium Processing Network



Network Input Task 1 processing nodes are loadable at run-time Task 2 each node defines its properties, parsed by driver Operation ID Code Task 1 Task 2 Operation driver schedules nodes as Function needed Critical Input Tasks Level high performance by efficient scheduling of computationally intensive nodes Processing Tracker Network Output

Processing Network

simple API for node functions

available as generic C library,

CPUs, e.g. LEON

API compatible to Xentium library

development on PC possible



int op_inc(unsigned long op_code, struct proc_task *t)

```
size t i:
                                                                                                               size_t n;
                                                                                                               unsigned int *p;
                                                             #define OP INC 0x12345678
                                                                                                               /* get the number of datums */
                                                             pt = pt_track_create(op_inc, OP_INC, 5);
                                                                                                               n = pt get nmemb(t);
                                                             pn_add_node(pn, pt);
                                                                                                               /* get the data buffer associated with this task */
                                                                                                               p = (unsigned int *) pt_get_data(t);
                                                                                                               /* n is not 0, but data is NULL, this task is malformed and
                                                                                                                * will be moved directly to the output node with all
                                                                                                                * elements set to zero
                                                                                                                */
                                                                                                               if (!p)
                                                                                                                      return PN_TASK_DESTROY;
                                                                                                               /* increment all items by 1 */
                                                                                                               for (i = 0; i < n; i++)
                                                                                                                      p[i]++;
can be used for processing on general-purpose
                                                                                                               /* signal successful completion */
                                                                                                               return PN_TASK_SUCCESS;
                                                                                                        }
                                                                                      struct proc_task *t = pt_create(data, 32, 4, 0, 0);
                                                                                      pt_add_step(t, OP_INC, NULL);
                                                                                      pt_add_step(t, OP_INC, NULL);
                                                                                      pt_add_step(t, OP_MUL, NULL);
                                                                                      pt_add_step(t, OP_DEC, NULL);
```

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pn_input_task(pn, t);



Outlook and Plans

- OS to be qualified and used with SMILE (joint ESA/CAS) Soft X-ray Imager (SXI) instrument's LEON3 DPU (MMU support, processing net)
- to be used in ATHENA (ESA-L) Wide Field Imager (WFI) Instrument Control and Power Unit (ICPU) with SSDP
- new development cycle started (SMILE features)
- follow-up development when actual SSDP hardware becomes available
- planned follow-up project: Xentium DSP function library



Questions?



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