## System Impact of Distributed Multicore Systems

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## Outline

- Study context
- Commercial embedded multi core processors
- Multi core for space use scenarios
  - Integrated payload controller / data processing
  - Extended IMA
  - Dynamic scheduling
- Xtratum port to NGMP
- NGMP Assessment



## **Study Context**

- System Impact of Distributed Multi core Systems objectives:
  - Multi core processor use scenarios
  - Port of an hypervisor on the NGMP
  - NGMP assessment

#### Industrial organization:

- Astrium (Prime): system analysis, hypervisor specification and use case implementation
- Universitat Politecnica de Valencia: hypervisor port to NGMP (Xtratum), hypervision technology expertise



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Study context

### Commercial embedded multi core processors

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# Existing commercial embedded multi core processors: state of the art

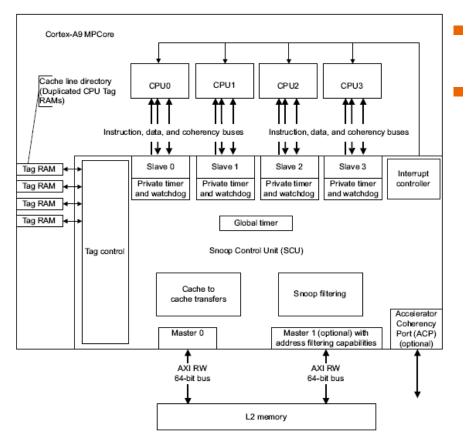
- Driven by mobile applications (smartphones, tablet, portable gaming console)
- Players:
  - CortexA9 MP (Apple A5, samsung exynos)
  - STEricsson Nomadik
  - Qualcomm Snapdragon

#### Commonalities:

- Homogeneous + specialized units
- ARM based
- Tightly coupled memory
- Multilayer AHB / AXI



## Existing commercial embedded multi core processors: example



Cortex-A9 MPCore

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# Existing commercial embedded multi core processors: programming model

#### SMP OS with APIs:

- PThread
- OpenMP
- MPI class of APIs

#### Parallelization principles

- Data parallelism: same code, different bits of data
- Task parallelism: different code, same data



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# Space domain needs for more capable processors

- Integrated Modular Avionics: centralization of processing tasks on a single computer
  - STR processing
  - GPS processing
- 100-1000 MIPS processor building block for payload processing
- NGMP is the solution currently being investigated



### Multi core use scenarios

#### Integrated payload controller / data processing

- Execution parallelism of controller and data processing
- Efficient implementation for small payload

#### Extended IMA application

- STR and GPS processing in the central OBC
- Partitioning of the CSW
- Stringent I/O requirements



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## Integrated payload controller / data processing: hypervisor

- Hypervisor (Xtratum) for T&SP partitioning of controller and data processing tasks:
  - Hypervisor allows using a guest OS for the controller and close to bare metal for processing
- Need for SMP partitions:
  - Typically, one core for controller, three cores for data processing



## Integrated payload controller / data processing: task parallelism

- Task parallelism is better suited than data parallelism for data processing on multi core
  - Explicit control of core execution: allows minimizing hardware concurrency
  - Smaller code fits in each core's L1 caches
  - Code rework to build SW pipeline



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## **Extended IMA application**

- Multiple partitions : STR, GPS, FDIR, AOCS, PAYLOAD ....
  - Hypervisor allows using different RTOS
  - Concurrent development and validation, IF HARDWARE COUPLING is low.
  - Indeterminism is dealt with by keeping huge performance margins.
- Processing requirements are low, constraints are on I/O latencies and throughput → use of DYNAMIC SCHEDULING



## Dynamic scheduling: I/O facts

#### I/Os are mostly not deterministic, eg:

- Flow controlled UARTs
- TM link bandwidth allocation mechanisms
- Asynchronous TC arrival

#### Therefore I/O operations are often IRQ driven



# Dynamic scheduling: time partitioning and IRQ handling on mono core

- Time partitioning  $\rightarrow$  IRQ handling latencies
  - In order not to break time partitioning, an asynchronous IRQ can only be processed during one of the execution slots of the destination partition
- IRQ handling latencies → performance loss or high hardware cost
  - To reduce IRQ handling latencies, it's possible to schedule partition that handles IRQ more often, but performances are degraded due to numerous context switches
  - Or I/Os must be more autonomous: large buffers, micro-controllers



# Dynamic scheduling: IRQ handling on multi core partitioned systems

- 1<sup>st</sup> possibility: I/O management partition running 100% on a dedicated core, problems:
  - Dedicating one complete core for IRQ results in performance loss of the overall system (especially on dual core)
  - IPC communications are needed to transfer data to the final user partition
  - Impossible to dedicate a given I/O to one partition, partitions are coupled via the I/O partition
- 2<sup>nd</sup> possibility: Dynamic scheduling on the core dedicated for IRQ handlers execution



## Dynamic scheduling: principles

#### As opposed to fixed cyclic scheduling:

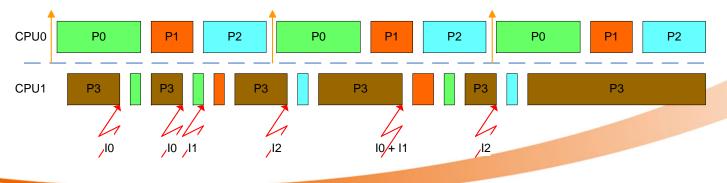
- Partitions scheduling is modified dynamically at runtime upon occurrence of external events (IRQ)
- Scheduling policy is chosen per core:
  - 3 cores using fixed cyclic scheduling, 1 core using dynamic scheduling



### Dynamic scheduling: use case

#### Partition execution is split

- Normal processing on cores using fix cyclic scheduling, IRQ handler execution on core using dynamic scheduling
- Background partition on core using dynamic scheduling
  - The dynamic scheduler executes the background partition when there isn't any IRQ handler to execute





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## Dynamic scheduling: benefits

- IRQ handling latency minimized
  - Allows lowering constraints on I/O architecture and scheduling plan

#### Efficient use of multi core processor

- A background partition can be executed between IRQs
- Efficient use of a additional core on strongly hardware coupled multi core processors

#### Efficient I/O operations

- No IPC needed, zero copy communication
- Allows for I/O direct allocation to partitions



## Dynamic scheduling: impacts

- Concurrent execution of IRQ handler and partition
  - Synchronization between IRQ handler task and other partition's tasks is achieved using hypervisor IPC or SMP OS primitives.

#### Time partitioning may be impacted

 The dynamic scheduler might not be able to enforce strict time partitioning, however the intrinsic hardware coupling of a multi core processor doesn't allow for strict time partitioning neither.



### Dynamic scheduling: other uses...

#### FDIR mechanisms

 Fixed cyclic scheduling plan doesn't allow for quick reaction to error conditions, dynamic scheduling can allow reducing fault mitigation latencies.

#### Enabler for advanced CSW partitioning



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### NGMP Assessment: IOMMU

#### IOMMU allows direct allocation of I/O to partitions

- Partitions can program DMA transfers in their virtual address space
- IOMMU ensures space partitioning
  - But not time partitioning! Rogue or failing partitions can "illuminate" the bus.



## NGMP Assessment: FLUSH instruction

- FLUSH = unprivileged instruction which flush both caches even when they are frozen
  - This breaks partitioning and forbids the possibility to prevent one partition to update the cache contents!
- SPARCv8 ISA doesn't require to flush the caches and moreover the FLUSH instruction is optional
- Some RTOS (RTEMS) are issuing flush instruction after installing trap handlers



## NGMP assessment: hardware coupling

- Hardware coupling: concurrency of access to the same resources
- Hardware coupling is an intrinsic property of multi core processors!
- Too much hardware coupling prevents concurrent development and validation
  - Partitions A and B are validated separately, if the hardware coupling is too important, at integration time the execution timings of A and B will change!



## NGMP assessment: hardware coupling

- Actual evaluation not yet performed
- Matters of concerns:
  - L1 caches size and write through policy: locality of space applications software is low
  - Shared AHB bus: round robin arbiter doesn't take into account the duration of the accesses
  - Single port L2 cache: the benefit of a cache hit for one core may be annulled by the fact it has to wait for a previous core to complete an external memory access.
  - Single memory controller: dual memory controller deemed necessary for simpler architecture like the SCOC3



## Conclusion

- Multi core execution parallelism is an enabler for some applications
- SMP partitions are needed to efficiently use all the cores
- IOMMU and dynamic scheduling should allow for efficient I/O handling mechanisms
- Hardware coupling can be a show stopper if too important



### Way forward

#### Evaluation of the performance of Xtratum port to the NGMP

 Use case implementation to benchmark I/O handling mechanisms



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