IXV DBSW -SOURCE CODE STATIC ANALYSIS





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IXVDBSW Source code STATIC ANALYSIS





- IXV OBSW and Spazio IT Code Quality Platform
- Bounded Model Checking and Abstract Intepretation (a proposed methodology)
- Analyses results
- Processes
- Spazio IT Code Quality Platforms @ AIRBUS Helicopters

Future Work

December 2014

IXV OBSW and Spazio IT Code Quality Platform





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IXV OBSW and Spazio IT Code Quality Platform



Spazio IT was requested to perform an activity of V&V on the entire IXV OVSW.

- To this purpose Spazio IT integrated the open source code quality platform SonarQube (<u>http://www.sonarqube.org</u>) with the following tools:
 - CppCheck (<u>http://cppcheck.sourceforge.net/</u>) open source – a C/C++ static analyser
 - PC-Lint (<u>http://www.gimpel.com/</u>) proprietary a rich pattern matching source code static analyzer (mostly used for MISRA C 2004 compliancy checks)

IXV OBSW and Spazio IT Code Quality Platform



- Spazio IT also integrated the following tools to see if they were applicable to the IXV OBSW and could provide additional information:
 - CBMC (<u>http://www.cprover.org/cbmc/</u>) open source – a C prover based on bounded model checking
 - Frama-C (<u>http://frama-c.com/</u>) open source a framework for the static analysis of C code – especially its "value analysis" (i.e. abstract interpretation) and "weakest precondition calculus" plugins.





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SonarQube – What is it?



SonarQube is an open source Web Application (<u>http://www.sonarqube.org</u>) which

- Takes in input a set of source code files and a set of analyses results (produced by external tools).
- Stores both sources and results in a database.
- Makes available the gathered information via a dynamic website where the results are shown in the context of the code itself.

SonarQube – What is it?





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SonarQube / Plugins / Sensors



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PC-Lint

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SonarQube – There's more



- Analyses on the same code base can be performed at different moments in time and SonarQube keeps track of the changes/evolution.
- The problems found during analyses (a.k.a. issues) can be managed directly from within the system itself, e.g.
 - Identifying false positives
 - Assigning issues to developers
 - Checking their status (if they have been solved)

- ...

http://sonarsrv.spazioit.com/



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http://sonarsrv.spazioit.com/

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- "Cppcheck is a static analysis tool for C/C++ code
- Unlike C/C++ compilers and many other analysis tools it does not detect syntax errors in the code.
- Cppcheck primarily detects the types of bugs that the compilers normally do not detect (from CppCheck website <u>http://cppcheck.sourceforge.net/</u>)"





- CppCheck produces messages belonging to 6 different categories:
 - 1. "Error used when bugs are found
 - 2. Warning suggestions about defensive programming to prevent bugs
 - 3. Style stylistic issues related to code clean-up (unused functions, redundant code, "constness", and such)
 - 4. Portability portability warnings. 64-bit portability. code might work different on different compilers. etc.
 - 5. Performance suggestions for making the code faster. These suggestions are only based on common knowledge.
 - 6. Information informational messages about checking problems."





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solution solution	nix.c portability nix.c portability hdog.c		835 scanf without field width limits can crash with huge input data of 890 scanf without field width limits can crash with huge input data of	n some
Summary: scanf v Message: scanf w %i => %3i Sample program t #include <stdio.h? int main() { int a; scanf("%i", &a) return 0; }</stdio.h? 	vithout field width limits can c ithout field width limits can cr that can crash: >	rash with huge input data on so ash with huge input data on libo	me versions of libc. : versions older than 2.13-25. Add a field width specifier to fix this problem:	-
To make it crash: perl -e 'print "5"x2	2100000' ./a.out			





- PC-Lint (<u>http://www.gimpel.com</u>) is a static analyzer deriving from the old Unix utility "lint".
- It supports many checks:
 - Barr Group's Top 10 Bug-Killing Rules (<u>http://www.barrgroup.com/webinars/10rules</u>)
 - Dan Saks (<u>http://www.dansaks.com/</u>)
 - Scott Meyers C++ books
 - (More) Effective C++ edition 1992
 - (More) Effective C++ edition 1996
 - (More) Effective C++ edition 1996

PC-Lint



- MISRA (http://www.misra.org.uk/)
 - Latest MISRA C (TM)
 - MISRA C 1998 (TM)
 - MISRA C 2004 (TM)
 - MISRA C 2012 (TM)
 - MISRA C++ 2008
 - MISRA C++ 2008 using 9000 level messages
- Barr's / Netrino Embedded C Coding Standard (<u>http://www.netrino.com/taxonomy/term/3</u>)
- Porting from 32-bit to 64-bit
 - from 32-bit to LP64 model
 - from 32-bit to LLP64 model
 - from 32-bit to ILP64 model





- In order to "reduce the noise", it needs a very careful setup / initial configuration (e.g.)
 - the proper memory model
 - the C/C++ include files (libraries)
 - the behaviour of standard macros and pragmas like "assert", "pragma pack",
 - the set of active checks

- ...

PC-Lint



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- "CBMC is a Bounded Model Checker for C and C++ programs. It supports C89, C99, most of C11 and most compiler extensions provided by GCC and Visual Studio. (...)
- It allows verifying array bounds (buffer overflows), pointer safety, exceptions and user-specified assertions.(..).
- The verification is performed by unwinding the loops in the program and passing the resulting equation to a decision procedure. (<u>http://www.cprover.org/cbmc/</u>)"





CBMC first converts a C program into a model, some kind of "symbolic executable".

Then the "symbolic executable" is "executed" and the execution generates "decision conditions" (about some property being true or false) that are expressed as CNF formulae.

Finally these formulae are passed to an external SAT solver for their evaluation and verification (<u>http://www.dwheeler.com/essays/minisat-user-guide.html</u>).





- Frama-C (<u>http://frama-c.com/</u>), like SonarQube, rather than being a specific tool, is "is an extensible and collaborative platform dedicated to source-code analysis of C software."
- Frama-C relies on CIL (C Intermediate Language) to generate an abstract syntax tree. The abstract syntax tree supports annotations written in ANSI/ISO C Specification Language (ACSL).





- Like SonarQube also Frama-C has its own Plugins. Spazio IT has used two Frama-C Plugins:
 - 1. Value analysis (<u>http://frama-c.com/value.html</u>) which computes a value or a set of possible values for each variable in a program. This plugin uses abstract interpretation techniques and many other plugins make use of its results.
 - WP (Weakest Precondition <u>http://frama-</u> <u>c.com/wp.html</u>) - to verify properties in a deductive manner

Methodology





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- Identify which checks need to be executed on the code, i.e.
 - for the compiler, which compiler warnings (possibly all of them) need to be verified;
 - for CppCheck, which type pf messages (errors, warnings, performance messages, and so on) need to be verified
 - for PC-Lint, which rule sets have to be used (e.g. MISRA C 2004), and for each rule set, which actual rules make sense and need to be verified





- Configure carefully the tools (in terms of tools options, selected memory model e.g. LP64 vs. LLP64, location of the sources, location of the include files, and so on...
- Tune/optimize the configuration identified in point 1 by running few analysis sessions to verify that the proper information is generated (and disable the production of useless, noisy outputs – this may require the development of some filtering scripts).
- Run the analyses whenever it makes sense in the lifetime of a project (or during operations), and possibly on a regular basis.





- At every run the code:
 - should compile;
 - should compile without generating any of the selected warning;
 - should pass CppCheck analyses without generating any of the selected messages;
 - should pass PC-Lint analyses without violating any of the selected rules/guidelines.

Basic Core



	Nov 13 2014 4.99.6-20141113-23:30	Nov 19 2014 4.99.6-20141119-11:30	Nov 21 2014 4.99.6-20141121-12:00	
Issues	15,763	15,394	15,306	
Blocker issues	0	0	0	(j <u></u>
Critical issues	144	0	0	~
Major issues	176	151	153	~
Minor issues	11,674	11,400	11,301	

 Gerard J. Holzmann , "Mars Code", Communications of the ACM, Vol. 57 No. 2, Pages 64-73, 10.1145/2560217.2560218 (http://cacm.acm.org/magazines/2014/2/171689-marscode/fulltext)

Bounded Model Checking Abstract Interpretation





When looking carefully into the magic ball...

what we eventually see...

is us. 🙂 🙂 🙂

Bounded Model Checking Abstract Interpretation



- CBMC and Frama-C Plugins (Value Analysis and WP) organize their computation into two phases:
 - Generation of a model of the code under analysis
 - "Symbolic execution" or "logic verification" of the model itself.
- The computation resources required by phase one grow in a polynomial way with the complexity of code under analysis (number of files, packages, classes, functions, parameters, variables, lines of code, loops, constructs and so o...)
- The computation resources required by phase two grow exponentially with the complexity of the code under of analysis.

Bounded Model Checking Abstract Interpretation



So, for not so small, real code bases

- either we stick to phase one
- or we split the system under analysis into reasonable, «manageable» chunks.



Never ending loop

```
#include <stdio.h>
```

```
int main() {
    int i = 0;
    int n = 10;
    for (i = 0; i < n; i++) {
        printf("Iteration #% 2d.\n", i + 1);
        if (i == 5) i = 0;
    }
    return 0;
}</pre>
```



■ CBMC reaction (phase one)

Unwinding loop c::main.0 iteration 1205 file loops.c line 7 function main thread 0 Unwinding loop c::main.0 iteration 1206 file loops.c line 7 function main thread 0 Unwinding loop c::main.0 iteration 1207 file loops.c line 7 function main thread 0 Unwinding loop c::main.0 iteration 1208 file loops.c line 7 function main thread 0 Unwinding loop c::main.0 iteration 1209 file loops.c line 7 function main thread 0

•••

■ Frama-C reaction (phase two)

•••

•••

- [value] Done for function printf
- [value] computing for function printf <- main. Called from loops.c:8.
- [value] Done for function printf
- [value] computing for function printf <- main. Called from loops.c:8.
- [value] Done for function printf
- [value] Recording results for main
- [value] done for function main
- [value] ====== VALUES COMPUTED ======
- [value] Values at end of function main:
 - NON TERMINATING FUNCTION

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■ Both CBMC and Frama-C would detect this

```
#include <stdio.h>
```

```
int main() {
  int i = 0;
  int n = 10;
  // int *pn = &n;
  int *pn;
  for (i = 0; i < (*pn); i++) {
    printf("Iteration #% 2d.\n", i + 1);
  }
  return 0;
}
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```

Loops Checking and Cyclic Tasks



If a CBMC analysis is conducted on a set of files using as entry point a given function (say 'foo') and the analysis finishes without having to limit neither the "unwinding" nor the "depth", then also the function 'foo' finishes, as well as any other function that is called by 'foo' (both directly or indirectly) and that belongs to the code under analysis.

Loops Checking and Cyclic Tasks



- Analysing the IXV OBSW it is possible to see that every cyclic task (cyclic thread) in the system is characterised by three functions:
 - 1. XXX_Init initialise the data required by the task
 - 2. XXX_StartThread start the thread
 - 3. XXX_Cycle this is the function that gets called at every cycle.
- So, if the CBMC analyses of all the XXX_Cycle functions finish, then (at least in the cyclic tasks) there is no neverending loop. And this is what was proved.

Manageable Chunks



■ Acting locally (at function level) → Identifying Manageable Chunks

■ Acting locally:

- pointer checks
- memory leak checks
- signed/unsigned overflow
- float overflow

- ...

Manageable Chunks





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Manageable Chunks



- The output produced by CBMC and Frama-C is similar to the one produced by static analysers based on pattern matching (like PC-Lint), similar but not the same, rather complementary.
- Frama-C Value Analysis Plugin did not bring interesting results when used at local level, at function level.





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Uninitialized Variables

- Compiler (gcc, clang, Visual Studio)
- PC-Lint

- Array Index out of bounds
 - PC-Lint in all code bases but only in simple cases
 - CBMC and Frama-C in all possible cases but in small portions of code



Constant Value Boolean Expression (MISRA C 2004 Rule 13.7)

```
void foo(unsigned int arg) {
...
if (arg < 0) {
    printf("Error: input parameter can only be a positive
integer\n");
    return;
  }
...
if (arg > UINT_MAX) ...
```

```
case SListIdx:
    tablePtr = &itPtr->sets;
    if (tablePtr != NULL) {
    if (arg > UINT_MAX) ...
```

```
– PC-Lint
```

Combining Signing and Unsigned Integers (MISRA C 2004 Rules 10.1, 10.3, 10.4)

- PC-Lint



- Implicit integer type conversion (and promotion) (MISRA C 2004 10.1, 10.3, 10.4, 10.6, 10.7, 10.8)
 - PC-Lint

Floating point comparison (MISRA C 2004 Rule 13.3)

double a; double b; if (a == b) ...

```
bool t My DoubleEquals (double64 t first, double64 t second)
ł
 bool t isEqual = FALSE;
  double64 t difference = My AbsValue(first - second);
  if (difference >= T EPSILON)
      isEqual = FALSE;
  else
      isEqual = TRUE;
  return isEqual;
}
    - PC-Lint
```

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PC-Lin

Problems with pointers

```
#include <stdio.h>
/* Why is this wrong? */
int main(void) {
    int x, *p;
    x = 10;
    *p = x;
    printf("*p = %d.\n", *p);
    return 0;
}
```

- PC-Lint
- CBMC / Frama-C



Divisions by Zero / Overflows

- PC-Lint
- CBMC / Frama-C
- Traps ←



Processes





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Who does what?



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All nowadays Integrated Development Environments (IDEs) like GNAT GPS 2014, Visual Studio 2013, Eclipse Luna, offer some form of Code Analysis.







IDE's analysis tools are to be used by software developers during their everyday work.

SonarQube analyses are more for the «quality people» and they are not supposed to be executed everyday, but rather at specific / well defined moments in the software development life cycle.





- SonarQube analyses should be performed after any «significant» delivery in a software development project, e.g. using ECSS 40 terminology, at:
 - CDR
 - QR
 - AR
- In maintenance projects SonarQube analyses should be performed after any «significant» new delivery, e.g. supposing a versioning like: *major.minor[.build[.revision]]* After every «minor» delivery.

AIRBUS Helicopters







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AIRBUS Helicopters



- Since mid 2012 Spazio IT has been working for AIRBUS Helicopters and has developed an Ada Plugin supporting both:
 - Adacore GNAT (<u>http://www.adacore.com</u>)
 - Atego APEX Ada (<u>http://www.atego.com</u>)
 compilation tools chains
- Spazio IT platform has been adopted by the group maintaining the software of the NH90 and Tiger helicopters.

Ada vs. C/C++





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- The majority if not all the problems mentioned earlier would never occur if Ada is used.
- In safety critical applications, the additional costs deriving by the adoption of Ada are partly compensated by the savings gained when performing Verification & Validation activities.
- What is the added value of using a code quality platform like the one developed by Spazio IT in the case of Ada?

■ The answer is: METRICS.





- Metrics like the lines of code, the % of comments, the cyclomatic complexity, the nesting and so on... are all correlated somehow to the readability and maintainability of the code.
- Being able to "see" these metrics in the context of the code allows developers and maintainers to immediately identify "host-spots", that is portions of code requiring attention.
- On top of that, the "time-machine" of SonarQube allows checking the evolution of these "hot-spots" with time.

Ada vs C/C++: origins and explanations



Real Time Execut Missile System User's Guide	ive for ms e	ten respective activity					
MC68020 C In	17 COSATI CODES			18. SUBJECT TERMS (Continue on reverse if necessary and identify by block number)			
Train Linearen	FIELD GROUP SUB-GROUP		RTEMS, real-ti	me, executive, heterogeneous, homogeneous,			
Two Press				multiprocessin	g, 68020, microprocessor, C language, runtime, (Continued on page ii)		
Relase 1.31 December 1991	Release 131 December 1991 Release 131 December 1991 Release 132 December 1991 Release 132 December 1991 Release 133 December 1991 Release 134 December 1991 Release 135 Release 135 December 1991 Release 135 December 1991 Release 135 Release 135 Release 135 Release 135 December 1991 Release 135 Release 135 Rele	rovides a ve, known itasking c riven, pri responsive ability. s of the d. The code en using the code we fine code we fine the code set of the set of the set of the the code we fine the set of the set of the set of the set of the set of the set of the set of the set of the set of the set of the se	high performant as RTEMS (Real apabilities; I ority-based, p e interrupt man RTEMS was or: Ada programmin de is Governme the 'C' progra as developed a lly RTEMS was	nce environment 1-Time Executive nomogeneous and preemptive sched anagement; dynam iginally develop ng language. RT ent owned, so no amming language as a linkable an developed for t	for embedded military applications. This for Missile Systems), includes such features heterogeneous multiprocessor systems; time uling; intertask communication and synchroni- ic memory allocation; and a high level of user ed in an effort to eliminate many of the major EMS is based on the RTEID (now ORKID) proposed licensing fees are necessary. The executive with a small amount of assembly language d/or ROMable library with the Ada programming he Motorola 68000 family of processors. It (Continued on page ii)		
		SSIFIED/UNLIMI	TED SAME AS F	PT DTIC USERS	UNCLASSIFIED		
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Ottobre 2013

Future Work





December 2014

Future Work



Continuous Integration

- Analyses should be run in an automatic way (e.g. via integration with systems like Jenkins - <u>http://jenkins-ci.org/</u>)
- More research on Abstract Interpretation
 - E.g. Clang Static Analyzer (<u>http://clang-analyzer.llvm.org/</u> a working example is available at http://sonarsrv.spazioit.com:8181/) vs. MathWorks Polyspace (<u>http://www.mathworks.com/products/polyspace/</u>)

Education

- C Awareness Campaign
- Code Quality Competence Centre
 - Training, Services, Platforms

Thank you for your time!



Software





December 2014

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